

Derived equivalences for derived discrete algebras of infinite global dimension

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 - if $X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} \Sigma X$ is exact in \mathcal{T} , then

$$F(X) \xrightarrow{F(f)} F(Y) \xrightarrow{F(g)} F(Z) \xrightarrow{\omega_X \circ F(h)} \Sigma' F(X)$$

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- $\eta: (F, \omega) \rightarrow (F', \omega')$ is a **natural transformation** between triangle functors if η is a natural transformation between F and F' st.
 $\omega' \circ \eta \Sigma = \Sigma' \eta \circ \omega$.

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Open question (Rickard '91)

Is any derived equivalence standard?

Affirmative answer for algebras:

- hereditary (Miyachi, Yekutieli '01),
- (anti-)Fano (Minamoto '12),
- triangular (Chen '16),
- **derived discrete of finite global dimension** (Chen, Zhang '19).

Definition (Vossieck '01)

A finite dimensional \mathbb{K} -algebra A is **derived discrete** if for any $\mathbf{n} = (n_i)_{i \in \mathbb{Z}}$, where $n_i \in \mathbb{N}_{\geq 0}$, there are only finitely many isoclasses of indecomposables in $\mathbf{D}^b(\text{mod } A)$ of cohomology dimension vector \mathbf{n} .

Example

The algebra of dual numbers given as a path algebra of the quiver:

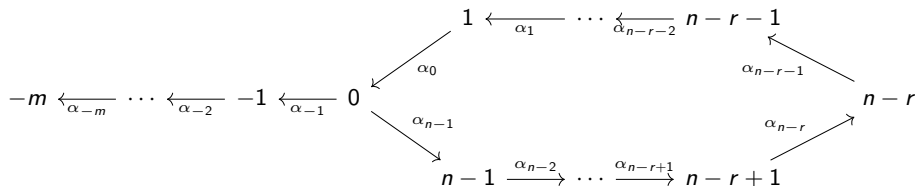
$$0 \rightrightarrows \alpha_0$$

where $\alpha_0^2 = 0$

Classification of derived discrete algebras

Let A be a finite dimensional algebra which is not piecewise hereditary of Dynkin type. Then A is derived discrete iff

$\mathbf{D}^b(\text{mod } A) \cong \mathbf{D}^b(\text{mod } \Lambda(r, n, m))$, where $1 \leq r \leq n$, $m \geq 0$ and $\Lambda(r, n, m) = \mathbb{K}\Delta(r, n, m)/I(r, n, m)$ is the path algebra of the quiver $\Delta(r, n, m)$:



bounded by $\alpha_{n-1}\alpha_0, \alpha_{n-2}\alpha_{n-1}, \dots, \alpha_{n-r}\alpha_{n-r+1}$.

Moreover, $\text{gldim } A = \infty$ iff $r = n$.

Ref: Vossieck '01 (Thm 2.1) and Bobiński-Geiss-Skowroński '04 (Thm A)

- Let \mathcal{A} be a category. A triangle functor $(F, \omega): \mathbf{K}^b(\mathcal{A}) \rightarrow \mathbf{K}^b(\mathcal{A})$ is a **pseudo-identity** if:
 - $F(X) = X$ for all objects $X \in \mathcal{A}$,
 - $F|_{\Sigma^n(\mathcal{A})} = \text{Id}_{\Sigma^n(\mathcal{A})}$.
- Note: any pseudo-identity is an autoequivalence.

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Theorem (Chen, Ye '18)

Let A be a finite dimensional algebra. Suppose that any pseudo-identity (F, ω) on $\mathbf{K}^b(\text{proj } A)$ is isomorphic to $(\text{Id}_{\mathbf{K}^b(\text{proj } A)}, \text{Id}_{\Sigma})$. Then for any finite dimensional algebra B any derived equivalence $\mathbf{D}^b(\text{mod } A) \xrightarrow{\sim} \mathbf{D}^b(\text{mod } B)$ is standard.

Strategy of the proof

Main result

Any derived equivalence between derived discrete algebras of infinite global dimension is standard.

Remark

A partial result has been obtained in [Chen, Ye, '18, Section 7] for the algebras of type $\Lambda(n, n, 0)$.

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We will start by describing a structure of $\text{ind } \mathbf{K}^b(\text{proj } \Lambda)$ for $\Lambda = \Lambda(n, n, m)$. Then:

- 0 pick an arbitrary pseudo-identity (F, ω) on $\mathbf{K}^b(\text{proj } \Lambda)$,
- 1 construct an isomorphism $\eta: F|_{\text{ind } \mathbf{K}^b(\text{proj } \Lambda)} \rightarrow \text{Id}_{\text{ind } \mathbf{K}^b(\text{proj } \Lambda)}$ (which extends to an isomorphism $\eta: F \rightarrow \text{Id}_{\mathbf{K}^b(\text{proj } \Lambda)}$),
- 2 then, by [Chen, Ye '18, Lemma 2.1] we obtain $(F, \omega) \xrightarrow{\sim} (\text{Id}, \omega')$ for some ω' ,
- 3 depending on n, m , either $\omega' = \text{Id}_\Sigma$ or $(\text{Id}, \omega') \cong (\text{Id}, \text{Id}_\Sigma)$.

Structure of the homotopy category

$\text{ind } \mathbf{K}^b(\text{proj } \Lambda) \cong \mathbb{K}\Gamma/\mathfrak{I}$ where $\Gamma = (\Gamma_0, \Gamma_1)$ and:

$$\Gamma_0 = \{(i, a, b) \mid i \in \{0, \dots, n-1\}, a \leq b + \delta_{i,0} \cdot m\}.$$

For every $v = (i, a, b) \in \Gamma_0$ we define two sets:

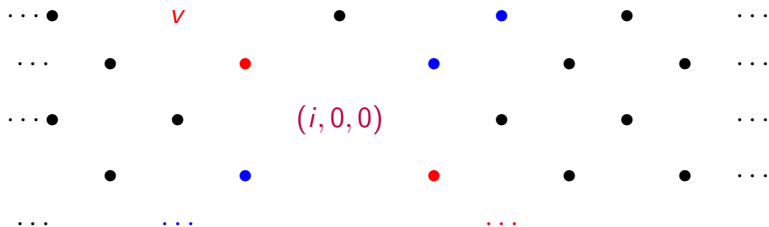
$$\mathcal{I}_v = \{(i, x, y) \mid x = a, \dots, b + \delta_{i,0} \cdot m, y \geq b\},$$

$$\mathcal{X}_v = \{(i+1, x, y) \mid x \leq a + \delta_{i,n-1} \cdot m, y = a, \dots, b + \delta_{i,0} \cdot m\}.$$

- for $u \in \mathcal{I}_v$ st. $u \neq v$ we have $f_{v,u}: v \rightarrow u$ in Γ_1 ,
- for $u \in \mathcal{X}_v$ we have $e_{v,u}: v \rightarrow u$ in Γ_1 ,
- note: $i+1$ is taken modulo n .
- **Ref:** Bobiński '11

Example \mathcal{I}_v and \mathcal{X}_v sets

$n = 1, m = 2, v = (0, 0, -m)$, we have sets \mathcal{I}_v and \mathcal{X}_v , intersection $\mathcal{I}_v \cap \mathcal{X}_v$ is in purple.



Structure of the homotopy category

The arrows are subject to the following relations for every $v \in \Gamma_0$:

- ① For $u \in \mathcal{I}_v, w \in \mathcal{I}_u, v \neq u, w \neq u$:

$$f_{u,w} \circ f_{v,u} = \begin{cases} f_{v,w} & w \in \mathcal{I}_v, \\ 0 & \text{otherwise.} \end{cases}$$

- ② For $u \in \mathcal{I}_v, w \in \mathcal{X}_u, v \neq u$:

$$e_{u,w} \circ f_{v,u} = \begin{cases} e_{v,w} & w \in \mathcal{X}_v, \\ 0 & \text{otherwise.} \end{cases}$$

- ③ For $u \in \mathcal{X}_v, w \in \mathcal{I}_u$:

$$f_{u,w} \circ e_{v,u} = \begin{cases} e_{v,w} & w \in \mathcal{X}_v, \\ 0 & \text{otherwise.} \end{cases}$$

- ④ For $u \in \mathcal{X}_v, w \in \mathcal{X}_u: e_{u,w} \circ e_{v,u} = 0$.

Structure of the homotopy category

- If $v = u$ then we define $f_{v,v}$ to be the identity morphism id_v .
- We write $e'_{v,u} = \begin{cases} e_{v,u} & u \in \mathcal{X}_v, \\ 0 & \text{otherwise.} \end{cases}$
- We have the following definitions of the suspension functor Σ and the AR-translation τ :
 - $\Sigma(i, a, b) = (i + 1, 1 + \delta_{i,n-1} \cdot m, 1 + \delta_{i,0} \cdot m)$
 - $\tau(i, a, b) = (i, a - 1, b - 1)$.

Lemma

Let $v = (i, a, b) \in \Gamma_0$. Then:

- 1 $f \in \text{End}(v)$ is an automorphism iff $f = \lambda f_{v,v} + \mu e'_{v,v}$ for $\lambda \neq 0$.
- 2 Suppose we have $f: v \rightarrow u$. Then f is irreducible iff $u \in \{(i, a+1, b), (i, a, b+1)\}$ and $f = \lambda f_{v,u} + \mu e'_{v,u}$ for $\lambda \neq 0$.

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Lemma

We have identification of projective modules up to Σ and AR-translation: P_{-j} with $(0, 0, -j)$ for $j = 0, \dots, m$ and P_k with $(k, 0, 0)$ for $k = 1, \dots, n-1$ (compare: Bobiński, '10, Corollary 6.3).

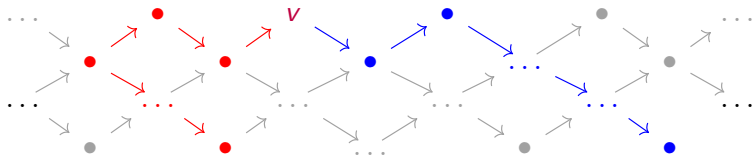
Proof – part 1

We construct $\phi_v \in \text{Aut}(v)$ for every $v \in \Gamma_0$, st. $F'(f) = f$ for every f where $F'(f) := \phi_u \circ F(f) \circ \phi_v^{-1}$ for $f: v \rightarrow u$.

Proof – part 1

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This will be done inductively. We put $\phi_v = \text{id}_v$ for $v = (i, 0, -\delta_{i,0} \cdot m)$. For the rest, we construct *canonical walks to and from* v for every connected component of the AR-quiver:



Fix $v = (i, a, b) \in \Gamma_0$ and let $f: v \rightarrow u$ be an irreducible morphism.

- 1 If $u = (i, a, b + 1) \in \Gamma_0$, then there exists $\phi \in \text{Aut}(u)$, st. $\phi \circ f = f_{v,u}$.
- 2 If $u = (i, a + 1, b + 1)$ for $b = a - \delta_{i,0} \cdot m$ such that $f \circ f_{w,v} = 0$ for $w = (i, a, b)$, then there exists $\phi \in \text{Aut}(u)$ st. $\phi \circ f = f_{v,u}$.

Dual statements for “backward” canonical walks also hold. For the proof of the above observation, let $f = \lambda f_{v,u} + \mu e'_{v,u}$, $\lambda \neq 0$. Then:

- 1 Put $\phi := \lambda^{-1} f_{u,u} - \lambda^{-2} \mu e'_{u,u}$.
- 2 Three subcases:
 - $e'_{u,u} = 0, e'_{v,u} = 0$. Proceed as in (1).
 - $e'_{u,u} = e_{u,u}, e'_{v,u} = e_{v,u}$. Proceed as in (1).
 - $e'_{u,u} = 0, e'_{v,u} = e_{v,u}$. This happens iff $n = 1, m = 1$. Note $u \notin \mathcal{I}_w$ but $u \in \mathcal{X}_w$. Hence $0 = f \circ f_{w,v} = \mu e_{w,u} \implies \mu = 0 \implies \phi := \lambda^{-1} f_{u,u}$.

By applying the above observations we inductively construct ϕ_u for every $u \in \Gamma_0$ such that $F'(f_{v,u})(= \phi_u \circ F(f_{v,u}) \circ \phi_v^{-1}) = f_{v,u}$.

If $n = 1$, then for brevity we will use $(a, b) := (0, a, b)$ for indecomposables.

In this step we will show that $F'(e_{v,u}) = e_{v,u}$ for every morphism $e_{v,u}: v \rightarrow u$.

General method

Find decomposition of morphisms of one of the forms:

- $e_{w_1,w_4} = f_{w_3,w_4} \circ e_{w_2,w_3} \circ f_{w_1,w_2}$ where either $e_{v,u} = e_{w_1,w_4}$ or $e_{v,u} = e_{w_2,w_3}$, or
- $f_{w_2,w_3} \circ e_{w_1,w_2} = e_{w,w_3} \circ f_{w_1,w}$ where either $e_{v,u} = e_{w_1,w_2}$ or $e_{v,u} = e_{w,w_3}$,

where, by induction, F' acts trivially on all the morphisms different from $e_{v,u}$. Obtain the result by applying F' to the both sides of the equation.

We now know that $(F, \omega) \cong (\text{Id}_{\mathbf{K}^b(\text{proj } \Lambda)}, \omega')$ for some ω' .

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Lemma

Let $(\text{Id}_{\mathbf{K}^b(\text{proj } \Lambda)}, \omega') : \mathbf{K}^b(\text{proj } \Lambda) \rightarrow \mathbf{K}^b(\text{proj } \Lambda)$. Then for every $v \in \text{ind}(\mathbf{K}^b(\text{proj } \Lambda))$ we get $\omega'_v = f_{\Sigma v, \Sigma v} + \mu e'_{\Sigma v, \Sigma v}$ where $\mu \in \mathbb{K}$. In particular if $n > 1$, then $\omega' = \text{id}_{\Sigma}$.

Idea of proof

$v = (i, a, b)$, $v = (i, a, b + 1)$, $w = (i, b + \delta_{i,0} \cdot m + 1, b + 1)$. There exists exact triangle: $v \xrightarrow{f_{v,u}} u \xrightarrow{f_{u,w}} w \xrightarrow{\lambda_{e_w, \Sigma v}} \Sigma v$ (see Bobiński, Schmude '20, Prop. 2.2). Apply $(\text{Id}_{\mathbf{K}^b(\text{proj } \Lambda)}, \omega)$ and use (TR3) axiom.

It suffices to consider $n = 1$.

- If $m > 0$, then in every ray of the AR-quiver we have $v = (a, b)$, st. $\dim_{\mathbb{K}}(\text{End}(v)) = 1$, hence $\omega'_v = \text{id}_{\Sigma v}$. If $f: v \rightarrow u$ for $u = (a, b + 1)$, then by functoriality $\omega'_u \circ \Sigma(f) = \Sigma(f)$. But since $\text{Hom}(v, u) \cong \text{End}(u)$, as left $\text{End}(u)$ -modules then $\omega'_u = \text{id}_{\Sigma u}$. Proceed by induction on b .

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- If $m = 0$, then $(\text{Id}_{\mathbb{K}^b(\text{proj } \Lambda)}, \omega') \cong (\text{Id}_{\mathbb{K}^b(\text{proj } \Lambda)}, \text{Id}_{\Sigma})$. Isomorphism is constructed iteratively. For $v = (0, b)$ $b \geq 0$ put $\eta_v := \text{id}_v$. Let $s \geq 0$. For $v = (s, b) \in \Gamma_0$ put:

$$\eta_v = \Sigma^{s-1}(\omega_{(0, b-s)}) \circ \Sigma^{s-2}(\omega_{(1, b-s+1)}) \circ \dots \circ \omega_{(s-1, b-1)}$$

For $v = (-s, b) \in \Gamma_0$ put:

$$\eta_v = \Sigma^{-s}(\omega_{(-1, b+s-1)}^{-1}) \circ \Sigma^{-s+1}(\omega_{(-2, b+s-2)}^{-1}) \circ \dots \circ \Sigma^{-1}(\omega_{(-s, b)}^{-1})$$